

# WIRELESS QUEUING NETWORKS UNDER CHURN AND MOTION

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## Structure of the Lecture

- Background and Motivation
- Methodology
- **Part 1**  
Birth and Death of Wireless Queues  
Joint work with A. Sankararaman and S. Foss
- **Part 2**  
Motion of Wireless Queues under Multihop Routing  
Joint work with S. Rybko, S. Shlosman and S. Vladimirov

## Motivations in Wireless Networks

- **Lack of understanding and analysis of**

  - Space-time interactions**

  - Static spatial setting well understood: Stochastic Geometry [FB, Blaszczyszyn 01]
  - Churn partly taken into account in flow-based queuing [Bonald, Proutiere 06], [Shakkottai, De Veciana 07]

- **Contents of this lecture:**

  - First models with such dynamics in stochastic geometry**

## Methodology

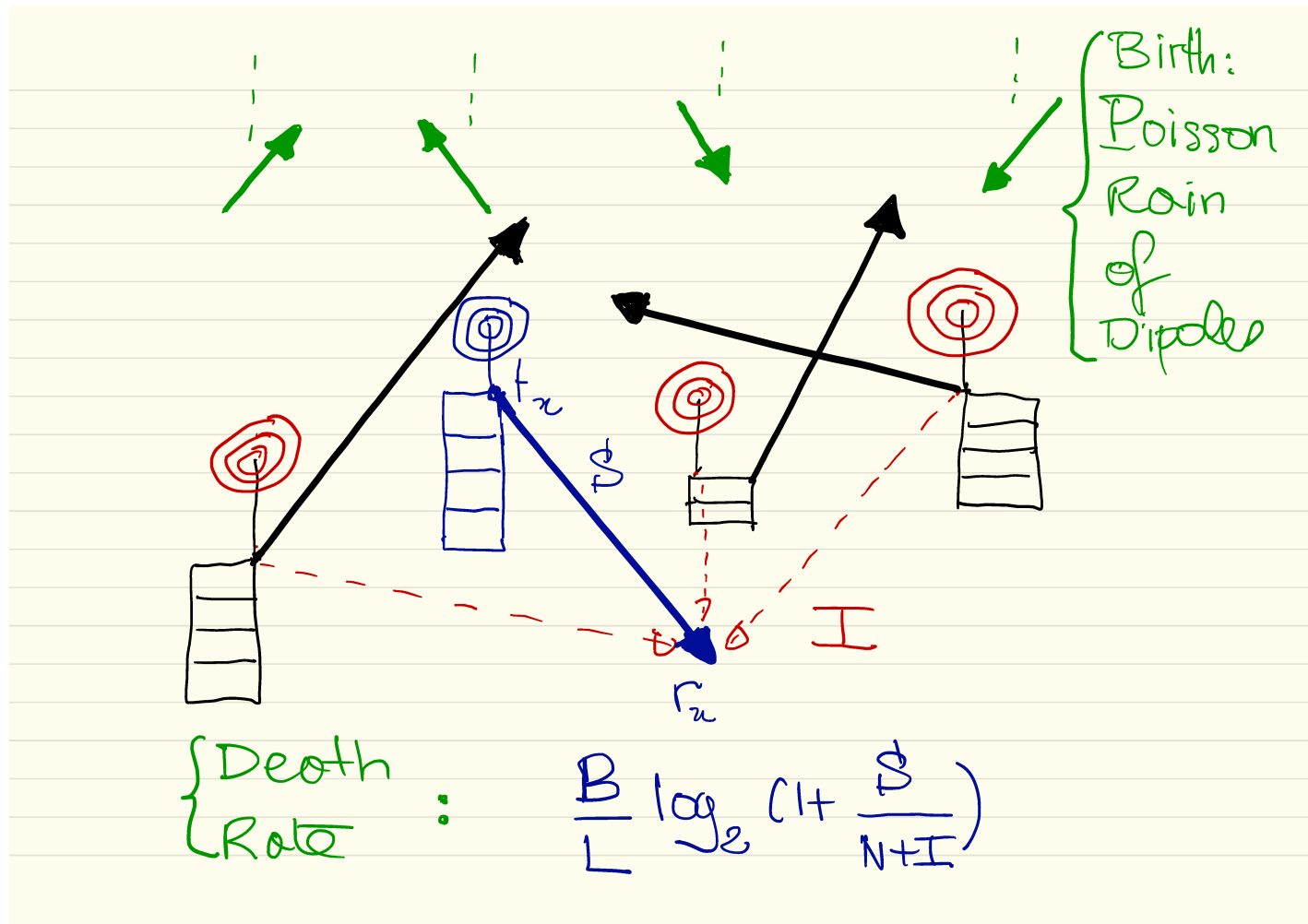
**Everything Should Be Made as Simple as Possible,  
But Not Simpler**  
**A. E.**

## Setting

- **Infrastructureless Wireless Network:**  
Ad-hoc Networks, D2D Networks, IoT
- **Markov Models:**  
Poisson, Exponential
- **Mathematical tools:**  
Stochastic Geometry, Fluid, Mean-Field

## I: Churn: Birth and Death

- Problem Statement
- Summary of Results
- Proof Overview



## Stochastic Network Model

- $S = [-Q, Q] \times [-Q, Q]$ : torus where the wireless links live
- **Links**: (Tx-Rx pairs)
- **Links**: **arrive** as a PPP on  $\mathbb{R} \times S$  with intensity  $\lambda$ :  
Prob. of a point arriving in space  $dx$  and time  $dt$ :  $\lambda dx dt$
- Each Tx has an **i.i.d. exponential file size**  
of mean **L** bits to transmit to its Rx
- A point **exits** after the Tx finishes transmitting its file
- $\Phi_t$ : set of locations of links present at time  $t$ :

$$\Phi_t = \{\mathbf{x}_1, \dots, \mathbf{x}_{N_t}\}, \quad \mathbf{x}_i \in S$$



## Interference and Service Rate

- Interference seen at point  $x$  due to configuration  $\Phi$

$$I(\mathbf{x}, \Phi) = \sum_{\mathbf{x}_i \in \Phi \neq \mathbf{x}} l(\|\mathbf{x} - \mathbf{x}_i\|)$$

- Distance on the torus
- $l(\cdot): \mathbb{R}^+ \rightarrow \mathbb{R}^+$ : path loss function

- The speed of file transfer by link at  $\mathbf{x}$  in configuration  $\Phi$  is

$$R(\mathbf{x}, \Phi) = B \log_2 \left( 1 + \frac{1}{N + I(\mathbf{x}, \Phi)} \right)$$

- $B, N$  Positive constants

## B& D Master Equation

- A point born at  $\mathbf{x}_p$  and time  $b_p$  with file-size  $L_p$  dies at time

$$d_p = \inf \left\{ t > b_p : \int_{u=b_p}^t \mathbf{R}(\mathbf{x}_p, \Phi_u) du \geq L_p \right\}$$

- **Spatial Birth-Death Process**

- Arrivals from the Poisson Rain
- Departures happen at file transfer completion

## Properties of the Dynamics

- The statistical assumptions imply that  $\Phi_t$  is a Markov Process on the set of simple counting measures on  $\mathcal{S}$
- **Euclidean extension** of the flow-level models of [Bonald, Proutiere 06], [Shakkottai, De Veciana 07]

## Questions

- **Existence and uniqueness** of the stationary regimes of  $\Phi_t$
- **Characterization** of the stationary regime(s) if existence

## Main Stability Results

$$a := \int_{\mathbf{x} \in S} l(\|\mathbf{x}\|) d\mathbf{x}$$

### ■ Theorem

- If  $\lambda > \frac{B}{\ln(2)La}$ , then  $\Phi_t$  admits no stationary regime.
- If  $\lambda < \frac{B}{\ln(2)La}$ , and  $r \rightarrow l(r)$  bounded and monotone, then  $\Phi_t$  admits a unique stationary regime

### ■ Sufficient condition by fluid limit

### ■ Corollary

For the path-loss model  $l(r) = r^{-\alpha}$ ,  $\alpha \geq 2$ , for all  $\lambda > 0$ , and all mean file sizes, the process  $\Phi_t$  admits no stationary-regime

## Main Qualitative Result

- $\Phi$  stationary point-process on  $S$  with Palm distribution  $\mathbb{P}^0$
- **Clustering**  
 $\Phi$  is clustered if for all bounded, positive, non-increasing functions  $f(\cdot) : \mathbb{R}^+ \rightarrow \mathbb{R}^+$ , the shot noise;

$$\mathbf{F}(\mathbf{x}, \Phi) := \sum_{\mathbf{y} \in \Phi \setminus \{\mathbf{x}\}} f(\|\mathbf{y} - \mathbf{x}\|)$$

satisfies

$$\mathbb{E}^0[\mathbf{F}(\mathbf{0}, \Phi)] \geq \mathbb{E}[\mathbf{F}(\mathbf{0}, \Phi)]$$

Main Qualitative Result (*continued*)

■ **Theorem**

The steady-state point process, when it exists, is **clustered**

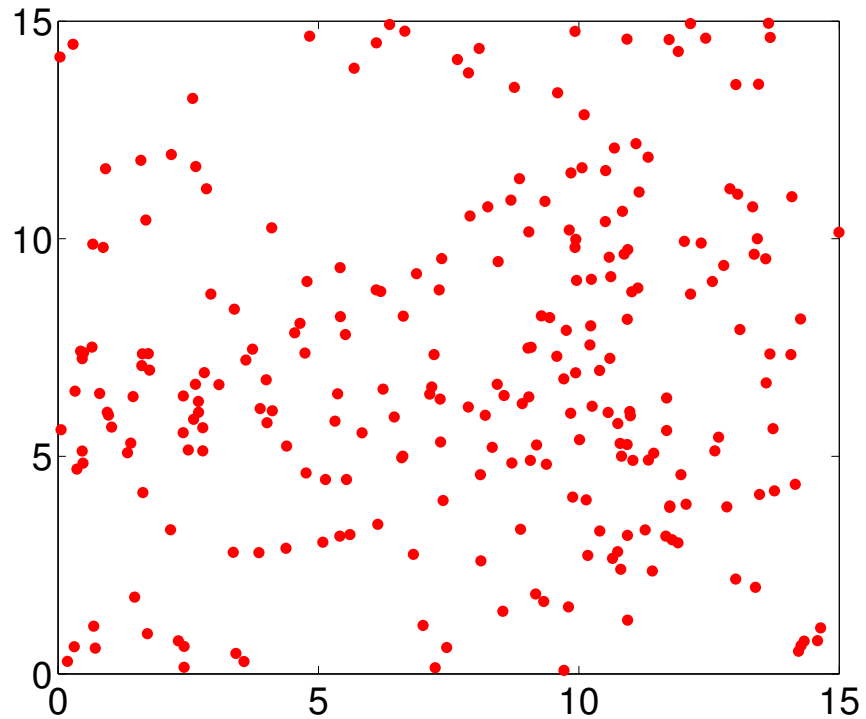
■ Follows from **Palm calculus + the FKG inequality**

■ **Interpretation of the result**

The steady-state interference measured at a uniformly randomly chosen point of is larger in mean than that at a uniformly random location of space.

■ **Key Observation**

- Dynamics Shapes Geometry
- Geometry Shapes Dynamics



A sample of  $\Phi$  when  $\lambda = 0.99$  and  $l(\mathbf{r}) = (\mathbf{r} + \mathbf{1})^{-4}$ .



## First Quantitative Results

- Mean-field approximations for the intensity of the steady-state process
  1. Poisson heuristic  $\beta_f$  - derived by neglecting clustering and assuming Poisson
  2. Second-order heuristic  $\beta_s$  based on a second-order cavity approximation of the dynamics

## Poisson Heuristic

**Exact Rate Conservation Law:**

$$\lambda \mathbf{L} = \beta \mathbb{E}_{\Phi}^0 \left[ \log_2 \left( 1 + \frac{1}{\mathbf{N} + \mathbf{I}(\mathbf{0})} \right) \right].$$

**Poisson Heur.:** Largest solution to the fixed point equation:

$$\lambda \mathbf{L} = \frac{\beta_f}{\ln(2)} \int_{z=0}^{\infty} \frac{e^{-Nz}(1 - e^{-z})}{z} e^{-\beta_f \int_{\mathbf{x} \in \mathcal{S}} (1 - e^{-zI(\|\mathbf{x}\|)}) d\mathbf{x}} dz$$

Ignores the Palm effect and uses that if  $X, Y$  are non-negative and independent,

$$\mathbb{E} \left[ \ln \left( 1 + \frac{X}{Y + a} \right) \right] = \int_{z=0}^{\infty} \frac{e^{-az}}{z} (1 - \mathbb{E}[e^{-zX}]) \mathbb{E}[e^{-zY}] dz.$$

## Second Order Heuristic

The intensity  $\beta_s$  is given by

$$\beta_s = \frac{\lambda L}{\mathbf{B} \log_2 \left( 1 + \frac{1}{\mathbf{N} + \mathbf{I}_s} \right)}$$

where  $\mathbf{I}_s$  is the smallest solution of the fixed-point equation

$$\mathbf{I}_s = \lambda L \int_{\mathbf{x} \in \mathcal{S}} \frac{\mathbf{I}(\|\mathbf{x}\|)}{\mathbf{B} \log_2 \left( 1 + \frac{1}{\mathbf{N} + \mathbf{I}_s + \mathbf{I}(\|\mathbf{x}\|)} \right)} d\mathbf{x}$$

## Second Order Heuristic (continued)

- **Rationale** based on  $\rho_2(\mathbf{x}, \mathbf{y})$ : second moment measure of  $\Phi$
- **Rate Conservation for  $\rho_2$** : when considering  $\mathbf{I}_s$  as a constant

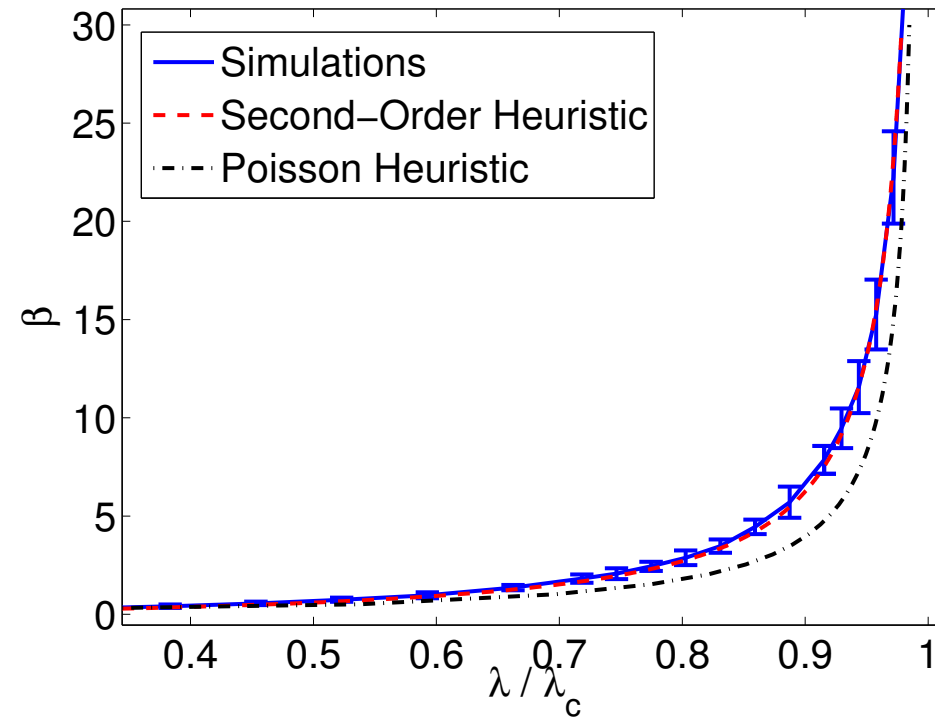
$$\rho_2(\mathbf{x}, \mathbf{y}) \frac{1}{L} \mathbf{B} \log_2 \left( 1 + \frac{1}{\mathbf{N} + \mathbf{I}_s + \mathbf{l}(\|\mathbf{x} - \mathbf{y}\|)} \right) = \lambda \beta_s$$

- From the definition of second moment measure,

$$\mathbf{I}_s = \int_{\mathbf{x} \in \mathbf{S}} \mathbf{l}(\|\mathbf{x}\|) \frac{\rho_2(\mathbf{0}, \mathbf{x})}{\beta_s} d\mathbf{x}$$

which gives the fixed point equation for  $\mathbf{I}_s$

- The formula for  $\beta_s$  follows from **Rate Conservation for  $\rho_1 = \beta_s$**



95% confidence interval when  $l(r) = (r + 1)^{-4}$

## Tightness Results & Extensions

- The Poisson heuristic is **tight** in heavy and light traffic
- **Recent Extensions** obtained with **S. Foss**:
  - **Exact expression** for the intensity  $\beta$  of  $\Phi$  in the **Low SINR** regime when replacing the death rate by

$$\frac{B}{\ln(2)L} \frac{S}{N + I(\mathbf{x}, \Phi)}$$

- **Scalability result**: extension to dynamics on  $\mathbb{R}^2$  using **Coupling from the Past** techniques.
- **Future**: introduction of **scheduling** or **multi-user IT**

## Summary

- A new basic representation of **space time interactions** in wireless networks
- A **generative model for clustering** as assumed in 3GPP simulation standards
- A new **dynamic notion of capacity** involving both queuing and IT
- **First analytical results** in the Low SINR case and good **heuristics** in general

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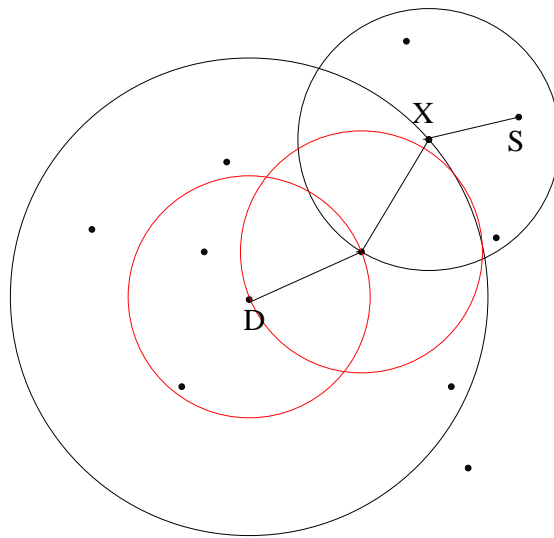
## II: Motion and Multi-Hop Routing

- **Problem Statement**
- **Summary of Results**

## Problem Statement

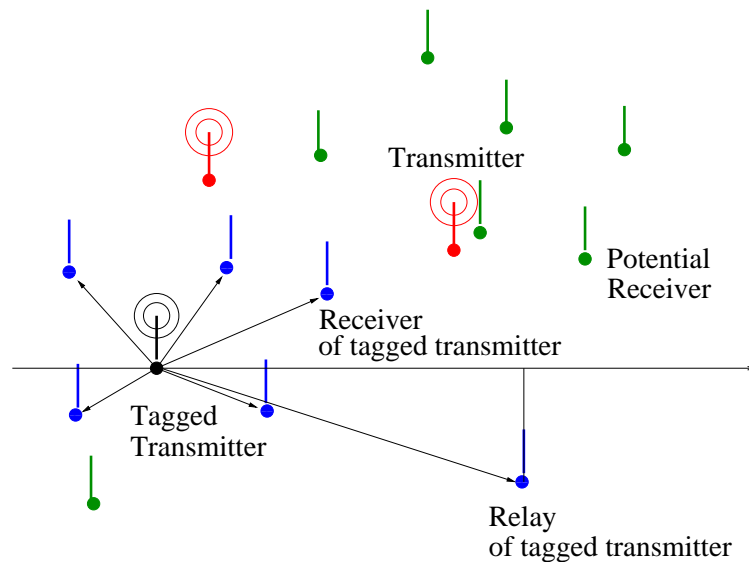
- **Setting:** Grossglauser & Tse 02 scaling law problem
  - Multihop relaying
  - Opportunistic geographic routing
  - Motion of nodes
- New **SG+QT** view of the problem

## Example of Geographic Routing



- **Nearest Neighbor Geographic Routing**  
The next hop on the route from  $S$  to  $D$  is the nearest among the nodes which are closer from  $D$  than  $X$ .
- On a Poisson P.P., a.s.
  - **No ties**
  - **Converges** in finite number of steps

## Wireless Geographic Routing



- Each node uses **Aloha** to split the Poisson p.p. into transmitters and potential receivers
- **Potential relays of a transmitter:** receivers with a large enough SINR
- **Geographic Routing:** next hop := potential relay nearest to destination

## Traffic and Relaying

- **DTN** type assumptions:
  - Wireless nodes **move**
  - Each moving node generates **packets**
  - Each generated packet has a **destination**, e.g. another node at finite distance
  - **Multihop relaying**: each node transmits
    - \* its own packets
    - \* those of other nodes on their route to destination
  - **Contention**: on each node, packets are queued **FIFO** and are served as above (SINR condition)

## Reduction (as per Methodology :)

- Wireless nodes move randomly on a grid or a graph  $G$  (e.g.  $\mathbb{Z}$  or  $\mathbb{Z}/K\mathbb{Z}$ ,  $\mathbb{Z}^2$ ,  $d$ -regular graph)
- **Traffic:**
  - Each moving node generates packets at rate  $\lambda$
  - Each generated packet has a destination (e.g. a point of the grid, vertex of the graph)
- **Contention:** on each node packets are queued **FIFO**

Reduction (as per Methodology :) ( <i>continued</i> )
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- **Motion:** neighbor nodes swap their positions with rate  $\beta$
- **Wireless:** communication to neighbors only
- **Opportunistic multihop routing:**  
upon service at a node, a packet:
  - is routed to the neighbor the **closest to its destination**
  - leaves the system if the destination is distance 1 or 0

## Finite Network Markovization

### ■ Assumptions

- Poisson arrivals with intensity  $\lambda$
- exponential service times with mean 1
- finite connected graph with  $K$  nodes

### ■ Markov representation with discrete non compact state:

- Permutation on  $[1, \dots, K]$   
(locations of wireless nodes)
- Ordered queue at each node  
(finite ordered list of destinations)



## Finite Network Instability Result

- Maximal degree in the graph:  $d$
- Motion rate  $\beta$
- **Theorem**

For all  $\beta > 0$ , this Markov process is transient when

$$K > K^* = \frac{d + 1}{\lambda}$$

- For all load factors, a finite network is unstable when the diameter of the network is large enough!

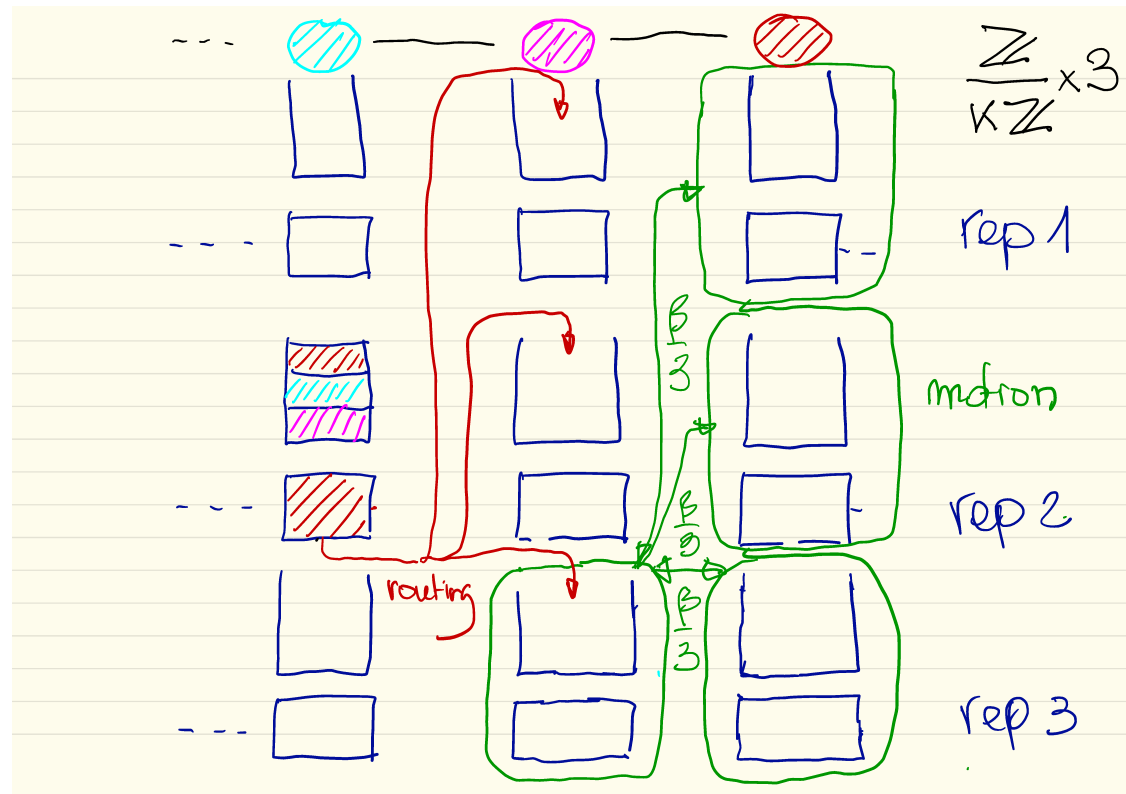
Finite Network Instability Result ( <i>continued</i> )
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■ Sketch of Proof in the  $d$ -regular case

- nodes mix to the uniform distribution on  $[1, \dots, K]$
- the proportion of time the server harboring a packet is a neighbor of its destination vertex is order  $\frac{d+1}{K}$ .
- if  $\lambda > \frac{d+1}{K}$ , drift inside orthant is positive on all coordinates
- **supermartingale** argument

## Replicated Version

- **N-Replica version** of the network on  $G$
- New graph  $G_N$  with
  - $V_N = V \times \{1, \dots, N\}$
  - edge between  $(u, i)$  and  $(v, j)$  if  $(u, v)$  edge in  $G$
- **Routing**
  - destination is any replica of the destination vertex in  $G$
  - routing to one of the  $N$  replicas closest to destination at random
- **Swapping**: between a replica and any neighboring replica at random



## The 3-Replica Version of $\mathbb{Z}/K\mathbb{Z}$

## Instability of the $N$ -Replica Version of a Finite Network

- Same instability result for  $G_N$  as for  $G$  for fixed  $N$
- **Example** The  $N$ -replica version of the network on  $\frac{\mathbb{Z}}{K\mathbb{Z}}$  is unstable under the same condition as the network on  $\frac{\mathbb{Z}}{K\mathbb{Z}}$  namely if  $K > 3/\lambda$

## Mean-Field Version

- **Mean-Field version** of the network on  $G$ :  
weak limit of the latter when  $N$  tends to infinity.
- **Notation:** network on  $G^\infty$
- **Existence:** tightness arguments

## Mean Field Networks on $\mathbb{Z}$

- **Non Linear Markov Process** roughly, dynamical system on probability measures  $\mu$  on queue states

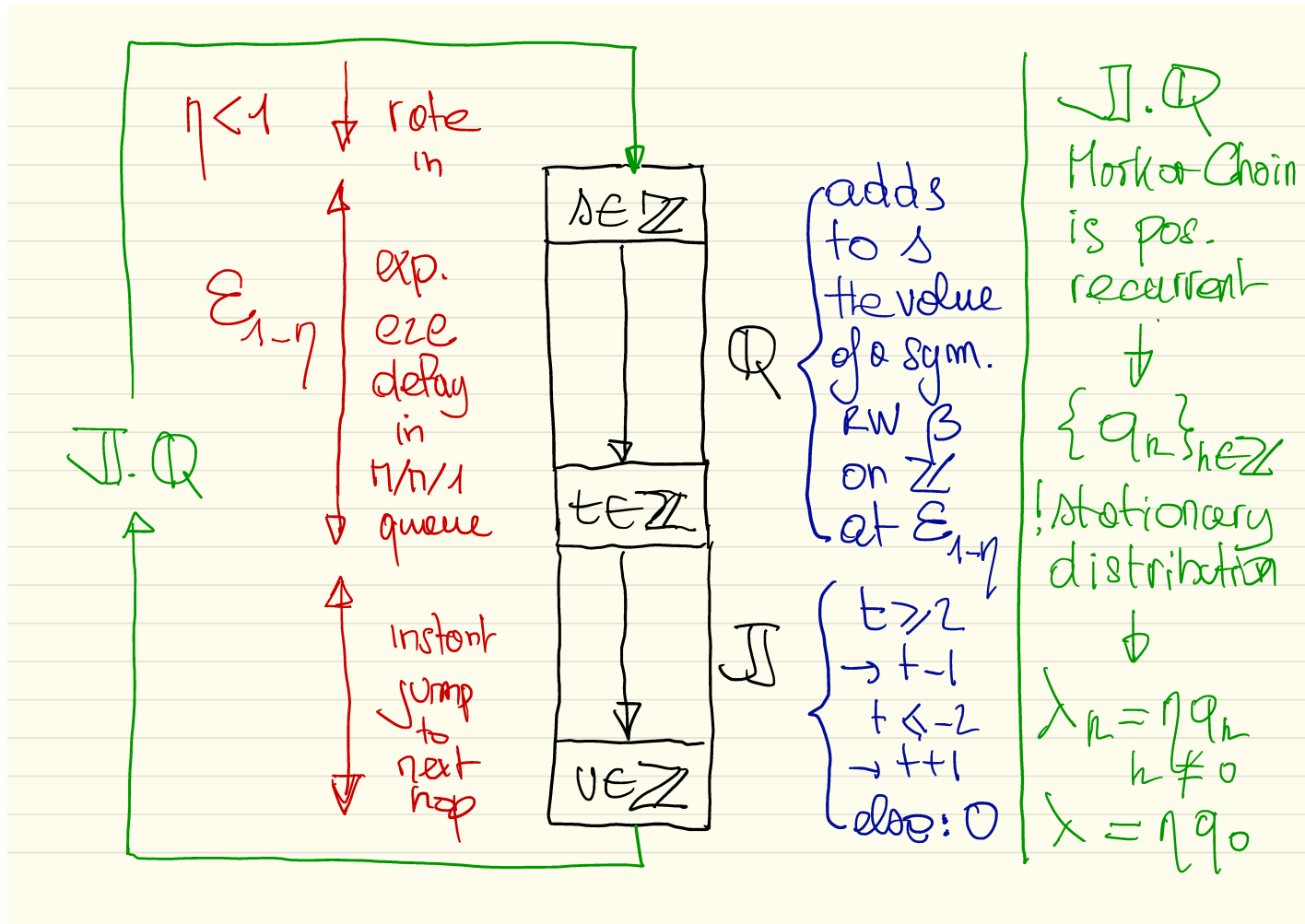
$$\mu(\mathbf{q}) = \mu(\mathbf{n}_1, \dots, \mathbf{n}_l), \quad \mathbf{n}_k : \text{relative location of dest}(\mathbf{c}_k)$$

- **Functional equation** for fixed points  $\mu(\mathbf{q})$  of this dynamical system

Mean Field Networks on  $\mathbb{Z}$  (continued)

- **Theorem** For all  $0 < \eta < 1$ , there exists
  - a unique  $0 < \lambda = \lambda_\eta < 1$
  - a unique probability distribution  $\mu = \mu_\eta$  on the queue state such that, for the exogenous arrival  $\lambda$ ,
    - $\mu$  is solution of the functional equation and
    - the total rate in a node under  $\mu$  is  $\eta$
- **Sketch of Proof**
  - Special case where the destination is the vertex of birth





## Existence of Multiple Solutions

- **Theorem** For the mean-field version of the network on  $\mathbb{Z}$ , there exists a  $\lambda_*$  such that for all  $\lambda < \lambda_*$ , there are at least two different values  $\eta = \eta_-(\lambda)$  and  $\eta = \eta_+(\lambda)$  s.t.

- $\lambda(\eta) = \lambda$

- $\eta_-(\lambda) \rightarrow 0$  as  $\lambda \rightarrow 0$

- $\eta_+(\lambda) \rightarrow 1$  as  $\lambda \rightarrow 0$

- **Sketch of Proof**

- When  $\eta$  tends to 0,  $\lambda(\eta) = \lambda q_0$  tends to 0 by M/M/1

- When  $\eta$  tends to 1,  $\lambda(\eta) = \lambda q_0$  tends to 0 by M/M/1 as well

## Illustration

- MC of the **distance to destination**:

$$\mathbf{p}(\mathbf{n}, \mathbf{n} + \mathbf{1}) = \beta, \quad \mathbf{p}(\mathbf{n}, \mathbf{n} - \mathbf{1}) = \beta + \gamma,$$

with  $\gamma = 1 - \eta$ ,

$$\mathbf{p}(\mathbf{1}, \mathbf{2}) = \beta, \quad \mathbf{p}(\mathbf{1}, \mathbf{0}) = \beta, \quad \mathbf{p}(\mathbf{1}, *) = \gamma$$

and

$$\mathbf{p}(\mathbf{0}, \mathbf{1}) = 2\beta, \quad \mathbf{p}(\mathbf{0}, *) = \gamma,$$

where  $*$  is absorbing

Illustration (continued)

- Mean # nodes visited by a customer till absorption:

$$\mathbb{E}[\mathbf{N}] = \mathbf{1} + \frac{\mathbf{2}\beta^2}{\gamma(\mathbf{3}\beta + \gamma)}$$

- Flow equation:

$$\eta = \mathbf{1} - \gamma = \lambda \left( \mathbf{1} + \frac{\mathbf{2}\beta^2}{\gamma(\mathbf{3}\beta + \gamma)} \right)$$

When  $\beta$  is large, there are two roots

$$\eta^- = \lambda + \frac{\mathbf{1} - \lambda - \sqrt{(\mathbf{1} - \lambda)^2 - \frac{\mathbf{8}}{\mathbf{3}}\lambda\beta}}{\mathbf{2}}$$

$$\eta^+ = \lambda + \frac{\mathbf{1} - \lambda + \sqrt{(\mathbf{1} - \lambda)^2 - \frac{\mathbf{8}}{\mathbf{3}}\lambda\beta}}{\mathbf{2}}$$

with

$$0 < \eta^- < \eta^+ < 1, \quad \nu^- \rightarrow 0, \quad \nu^+ \rightarrow 1, \quad \text{when } \lambda \rightarrow 0$$

## Generalization

- The multiple fixed point result can be extended to the Cayley graph of any group  $G$  s.t.
  - $G$  has a finite generating set
  - $G$  is infinite
  - the arrival rate, swap rate, swap rule, destination rule are  $G$ -invariant
- **Example:** the network of  $(\mathbb{Z})^\infty$  has at least two fixed points stationary regimes for every  $\lambda < \lambda_*$  with  $\lambda_* > 0$ .

## Meta-Stability

- Finitely many replicas–Infinitely many replicas stability difference.
- No contradiction with the fact that, for  $N < \infty$ , the network of  $(\mathbb{Z})^N$  has no stationary regimes  
**the time - replica diagram does not commute here!**

## Current and Future Extensions

- **Wireless primitives:** (SINR, MAC) to assess transmission between nearby nodes
- **Motion primitives:** Brownian, random waypoint
- **Queuing primitives:** beyond Poisson arrivals and exponential service times
- **Scheduling:** beyond FIFO: needed to see why motion increases capacity

## Summary

- **Way to go** but first step of a mean-field representation of **mutihop routing** in wireless networks beyond scaling laws
- **Meta-stability** is encouraging news
- When fully interconnected with SG, new **analytical handle** for optimization in this class of problems.



## References

### ■ Part 1

- with **A. Sankararaman**,  
Spatial birth-death wireless networks,  
**ArXiv 1604.07884**, under revision for **IEEE Tr. IT**
- with **S. Foss and A. Sankararaman**,  
Infinite spatial birth-death wireless networks,  
in preparation

### ■ Part 2

- with **S. Rybko, S. Shlosman and S. Vladimirov**,  
Metastability of Queuing Networks with Mobile Servers,  
**ArXiv 1704.02521**, submitted